

# Fuzzing Low-Level Code



**EPFL**



hexhive

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<https://hexhive.github.io>

# HexHive is hiring!



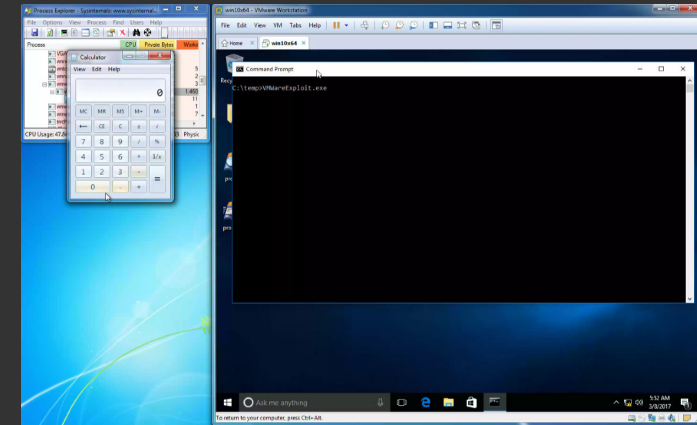
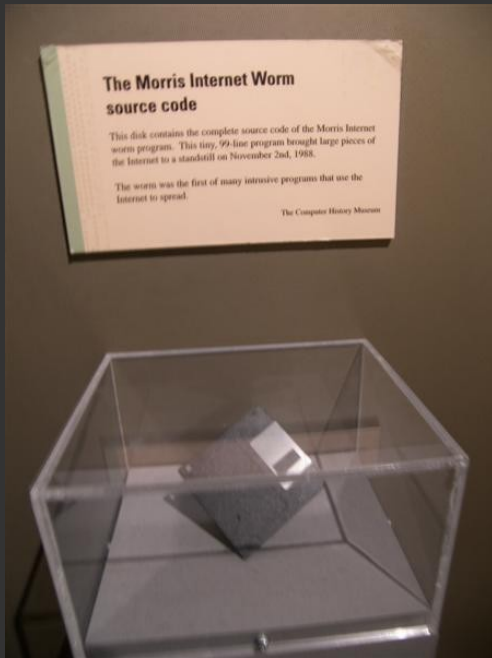
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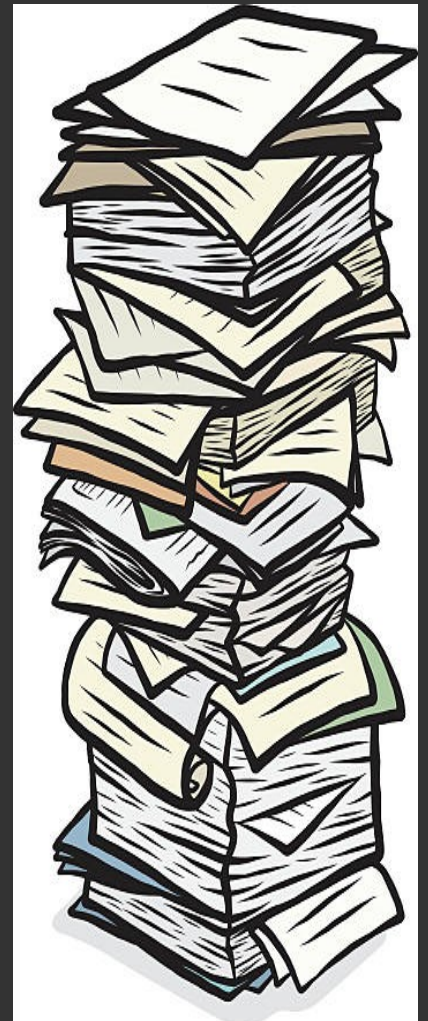
# Challenge: vulnerabilities everywhere



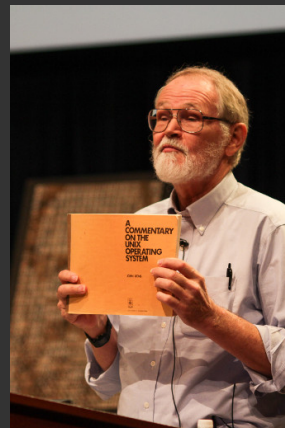
# Challenge: software complexity

<b>Google Chrome:</b>	<b>76 MLoC</b>
<b>Gnome:</b>	<b>9 MLoC</b>
<b>Xorg:</b>	<b>1 MLoC</b>
<b>glibc:</b>	<b>2 MLoC</b>
<b>Linux kernel:</b>	<b>17 MLoC</b>

Chrome and OS  
~100 mLoC,  
27 lines/page,  
0.1mm/page  $\approx$  370m

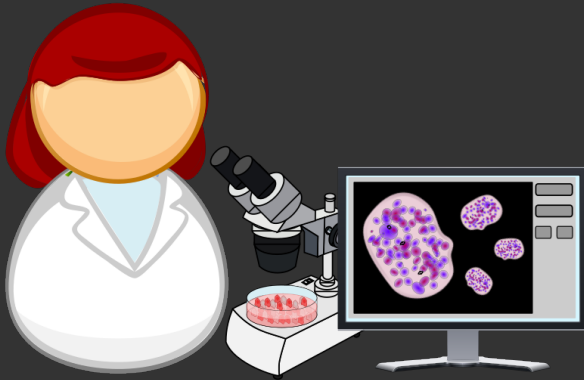


Margaret Hamilton  
with code for Apollo  
Guidance Computer  
(NASA, '69)



Brian Kernighan holding  
Lion's commentary on  
BSD 6 (Bell Labs, '77)

# Defense: Testing *OR* Mitigating?



## Software Testing

```
vuIn("AAA");
```

```
vuIn("ABC");
```

```
vuIn("AAAABBBB");
```

```
strcpy_chk(buf, 4, str);
```



## Mitigations

```
C/C++  
void log(int a) {  
    printf("A: %d", a);  
}  
  
void vuIn(char *str) {  
    char *buf[4];  
    void (*fun)(int) = &log;  
    strcpy(buf, str);  
  
    fun(15);  
}
```

```
CHECK(fun, tgtSet);
```

# Status of deployed defenses

- Data Execution Prevention (DEP)
- Address Space Layout Randomization (ASLR)
- Stack canaries
- Safe exception handlers
- Control-Flow Integrity (CFI):  
Guard indirect control-flow

## Memory

0x400

R-X

text

0x800

RW~~X~~

data

0xf??

RW~~X~~

stack

# Assessing exploitability





# Which crash to focus on first?

american fuzzy lop 2.32b (test_decode_bmp)			
process timing		overall results	
run time : 0 days, 1 hrs, 53 min, 36 sec		cycles done : 2	
last new path : 0 days, 0 hrs, 0 min, 35 sec		total paths : 939	
last uniq crash : 0 days, 0 hrs, 6 min, 18 sec		uniq crashes : 124	
last uniq hang : 0 days, 0 hrs, 16 min, 41 sec		uniq hangs : 128	
cycle progress		map coverage	
now processing : 120* (12.78%)		map density : 0.23% / 1.45%	
paths timed out : 0 (0.00%)		count coverage : 4.75 bits/tuple	
stage progress		findings in depth	
now trying : bitflip 1/1		favored paths : 126 (13.42%)	
stage execs : 923/5152 (17.92%)		new edges on : 185 (19.70%)	
total execs : 11.2M		total crashes : 7089 (124 unique)	
exec speed : 3487/sec		total hangs : 68.3k (128 unique)	
fuzzing strategy yields		path geometry	
bit flips : 291/1.27M, 56/1.26M, 22/1.26M		levels : 11	
byte flips : 7/158k, 16/29.9k, 23/30.3k		pending : 644	
arithmetics : 100/1.66M, 8/1.59M, 99/1.18M		pend fav : 0	
known ints : 4/93.8k, 22/395k, 61/768k		own finds : 938	
dictionary : 0/0, 0/0, 0/0		imported : n/a	
havoc : 353/1.43M, 0/0		stability : 100.00%	
trim : 19.01%/76.7k, 80.64%			
^C		[cpu000: 50%]	



# Residual Attack Surface Probing

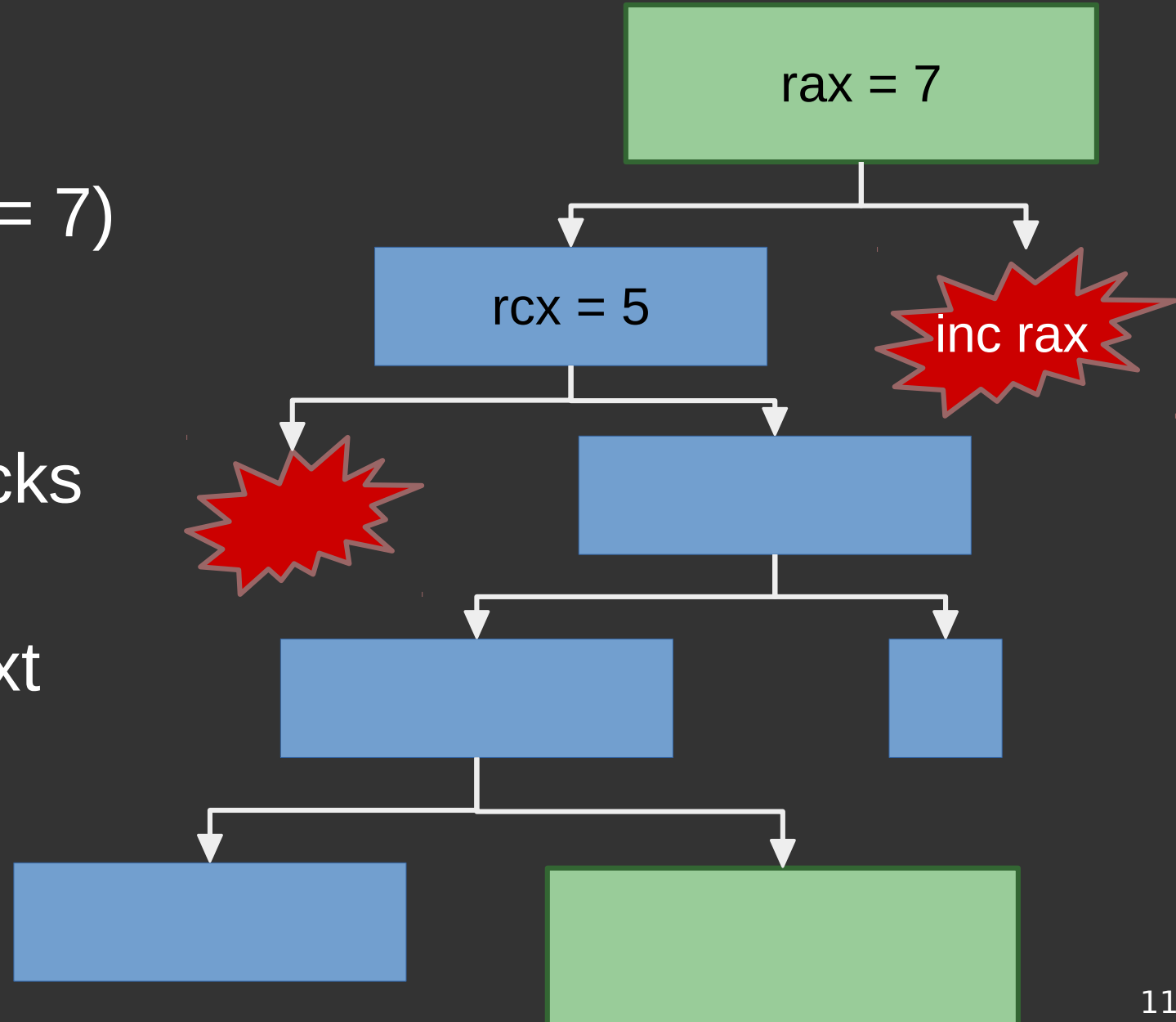
- State-of-the-art mitigations complicate attacks
  - Mitigations have limitations but these are hard to assess and explore systematically (and globally)
- Let's infer the ***Residual Attack Surface***
  - Given a crash/bug what can an adversary still do?
  - Residual attack surface depends on program, environment, and input

# Approach in a nutshell

- Given: crash that results in arbitrary write
- Assume: mitigations make exploitation hard
- Perform ***Code Reuse*** using ***Data-Only Attack***
  - Leverage memory corruption to corrupt state
  - Build Turing-complete payloads as execution traces
  - Express execution traces as memory writes

# BOP Gadget: basic block sequence

- **Functional:**  
compute (rax = 7)
- **Dispatcher:**  
connect  
functional blocks
- **Clobbering:**  
destroy context



***SPL  
payload***



```
graph TD; A["SPL payload"] --> B["Selecting functional blocks"]; B --> C["Searching for dispatcher blocks"]; C --> D["Stitching BOP gadgets"]
```

The diagram illustrates a four-step process for handling an SPL payload. It begins with a box labeled 'SPL payload' in the top-left corner, which has a yellow border. An arrow points from this box to a box labeled 'Selecting functional blocks' in the top-right corner. From there, an arrow points down to a box labeled 'Searching for dispatcher blocks' in the bottom-right corner. Finally, an arrow points from this box to a box labeled 'Stitching BOP gadgets' in the bottom-left corner. All boxes are rounded rectangles with black borders, except for the first one which has a yellow border.

Selecting  
functional blocks

Stitching  
BOP gadgets

Searching for  
dispatcher blocks

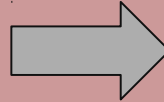


# SPL payload

- Payload language
- Subset of C
- Library Calls
- Abstract registers as volatile vars

```
void payload() {  
    string prog = "/bin/sh\0";  
    int64* argv = {&prog, 0x0};  
  
    __r0 = &prog;  
    __r1 = &argv;  
    __r2 = 0;  
  
    execve(__r0, __r1, __r2);  
}
```

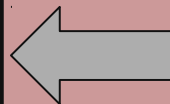
***SPL  
payload***



***Selecting  
functional blocks***



**Stitching  
BOP gadgets**



**Searching for  
dispatcher blocks**

# Functional block selection

- Find set of candidate blocks for SPL statement
- Candidate blocks “***could be***” functional blocks as they execute the correct computation
- What about other side effects? What about chaining functional blocks?

# Functional block selection (example)

     **r0** = 10;  
     **r1** = 20;

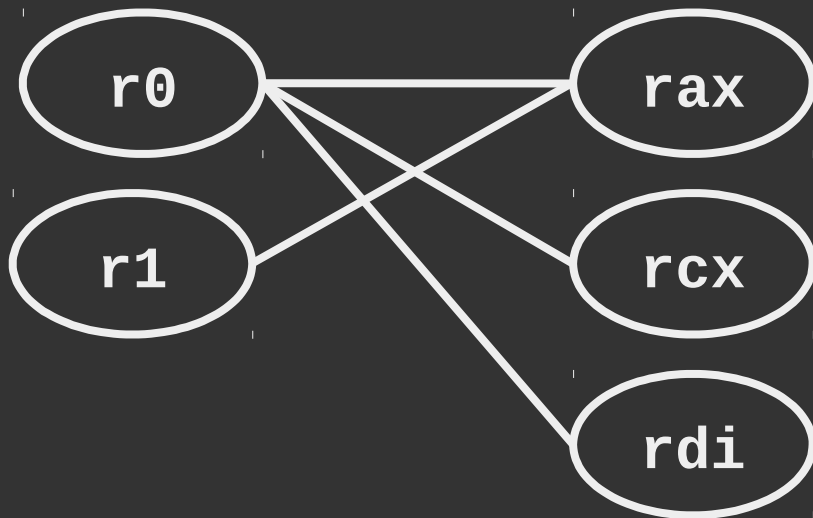
rax = 10

rdi = 10

rax = 20

rcx = 10

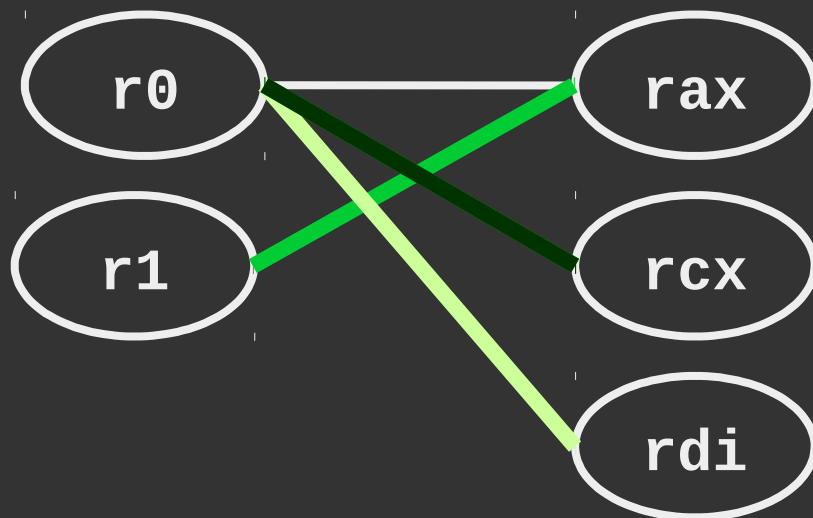
rcx = 30





# Functional block selection (example)

```
__r0 = 10;  
__r1 = 20;
```



`rax = 10`

Clobbering Clobbering

`rdi = 10`

Dispatcher Functional

`rax = 20`

Functional Functional

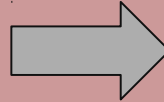
`rcx = 10`

Functional Dispatcher

`rcx = 30`

Clobbering Dispatcher

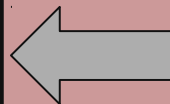
***SPL  
payload***



***Selecting  
functional blocks***



**Stitching  
BOP gadgets**

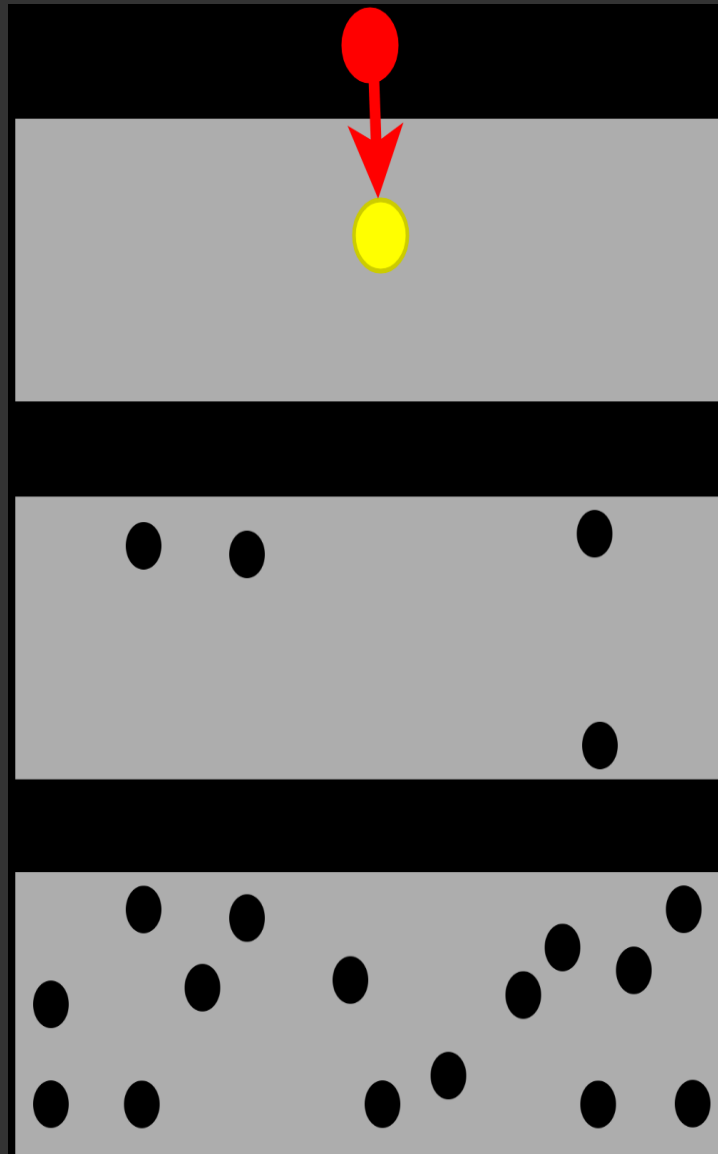


***Searching for  
dispatcher blocks***

# Dispatcher block search

- BOP gadgets are *brittle*
- Side-effects make gadgets hard to chain
  - Stitching gadgets is NP-hard
  - There is no approximative solution
- Our approach: back tracking and heuristics

# BOP gadgets are brittle



**Statement #1**

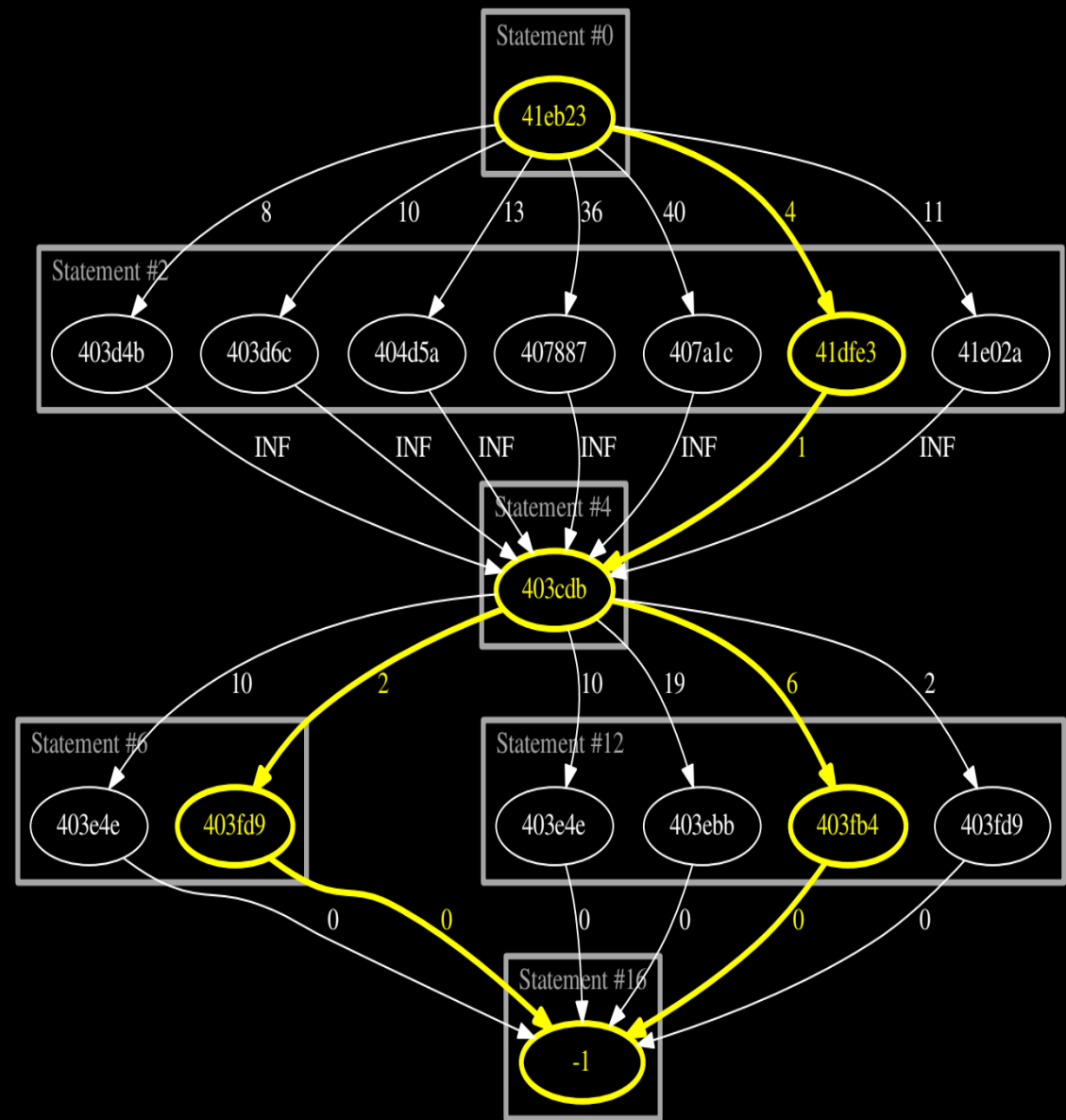
**Statement #2**

**Statement #3**

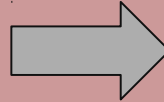


# Delta Graph: keeping track of blocks

- Squares: Functional blocks for SPL statements
- Nodes: Functional blocks
- Edges: Length of dispatcher chain
- Goal: Select one “node” from each layer (yellow)



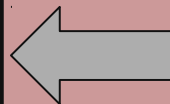
***SPL  
payload***



***Selecting  
functional blocks***



***Stitching  
BOP gadgets***



***Searching for  
dispatcher blocks***

# Stitching BOP gadgets

- Each path is a candidate exploit
- Check and validate constraints along paths
  - Goal: find a valid configuration
  - Constraints come from environment, SPL program, or execution context
  - Verify using concolic execution & constraint solving

# Payload synthesis

Program	SPL payload												
	<i>regset4</i>	<i>regref4</i>	<i>regset5</i>	<i>regref5</i>	<i>regmod</i>	<i>memrd</i>	<i>memwr</i>	<i>print</i>	<i>execve</i>	<i>abloop</i>	<i>infloop</i>	<i>ifelse</i>	<i>loop</i>
ProFTPD	✓	✓	✓	✓	✓	✓	✓	✓ 32	X <sub>1</sub>	✓ 128+	✓ ∞	✓	✓ 3
nginx	✓	✓	✓	✓	✓	✓	✓	X <sub>4</sub>	✓	✓ 128+	✓ ∞	✓	✓ 128
sudo	✓	✓	✓	✓	✓	✓	✓	✓	✓	X <sub>4</sub>	✓ 128+	X <sub>4</sub>	X <sub>4</sub>
orzhttpd	✓	✓	✓	✓	✓	✓	✓	X <sub>4</sub>	X <sub>1</sub>	X <sub>4</sub>	✓ 128+	X <sub>4</sub>	X <sub>3</sub>
wuftdp	✓	✓	✓	✓	✓	✓	✓	✓	X <sub>1</sub>	✓ 128+	✓ 128+	X <sub>4</sub>	X <sub>3</sub>
nullhttpd	✓	✓	✓	✓	✓	✓	X <sub>3</sub>	X <sub>3</sub>	✓	✓ 30	✓ ∞	X <sub>4</sub>	X <sub>3</sub>
opensshd	✓	✓	✓	✓	✓	✓	X <sub>4</sub>	X <sub>4</sub>	X <sub>4</sub>	✓ 512	✓ 128+	✓	✓ 99
wireshark	✓	✓	✓	✓	✓	✓	✓	✓ 4	X <sub>1</sub>	✓ 128+	✓ 7	✓	✓ 8
apache	✓	✓	✓	✓	✓	✓	✓	X <sub>4</sub>	X <sub>4</sub>	✓ ∞	✓ 128+	✓	X <sub>4</sub>
smbclient	✓	✓	✓	✓	✓	✓	✓	✓ 1	X <sub>1</sub>	✓ 1057	✓ 128+	✓	✓ 256

✓ The SPL payload was successfully executed on the target binary

X<sub>1</sub> Not enough candidate blocks

X<sub>2</sub> No valid register/variable mappings

X<sub>3</sub> No valid paths between functional blocks

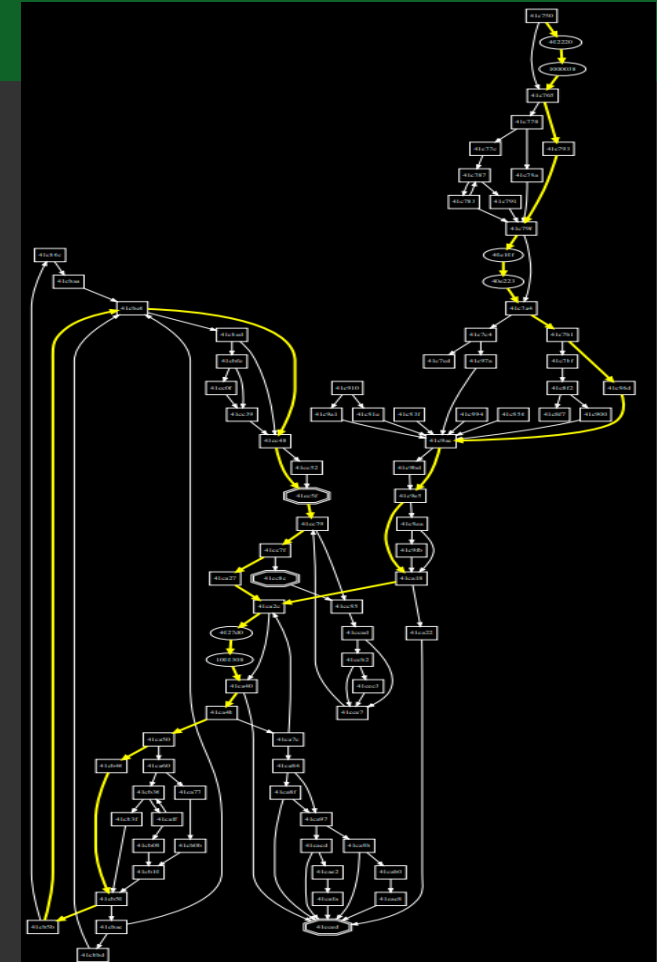
X<sub>4</sub> Un-satisfiable constraints or solver timeout

**Success Rate: 81%**

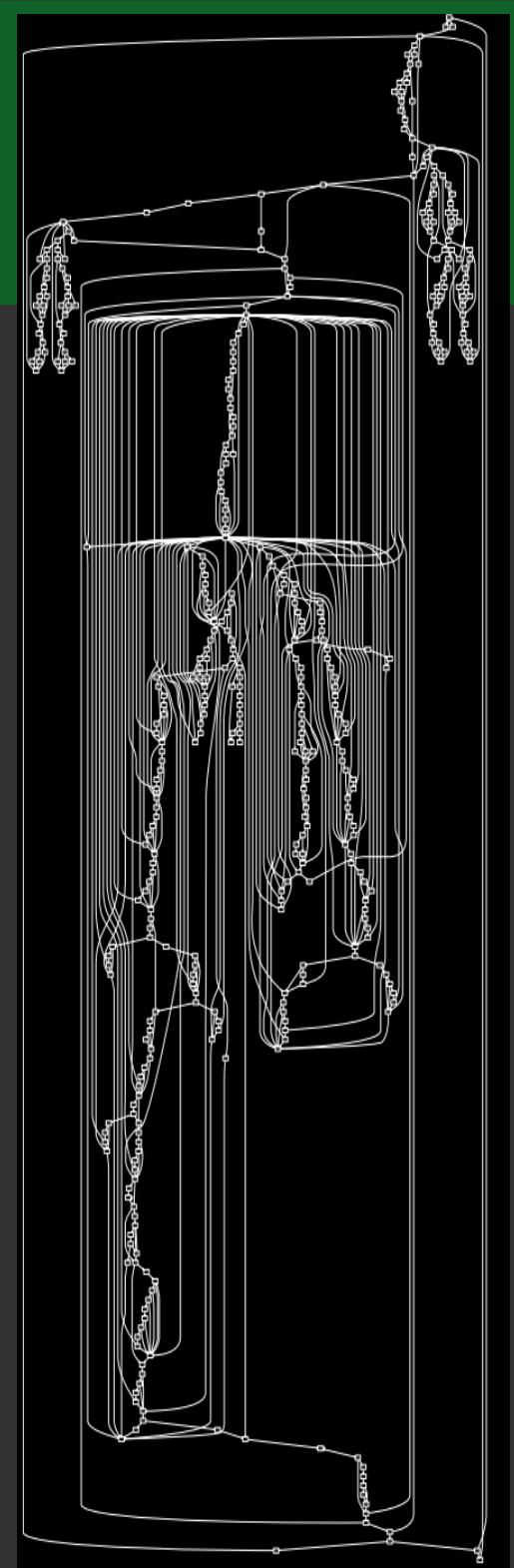
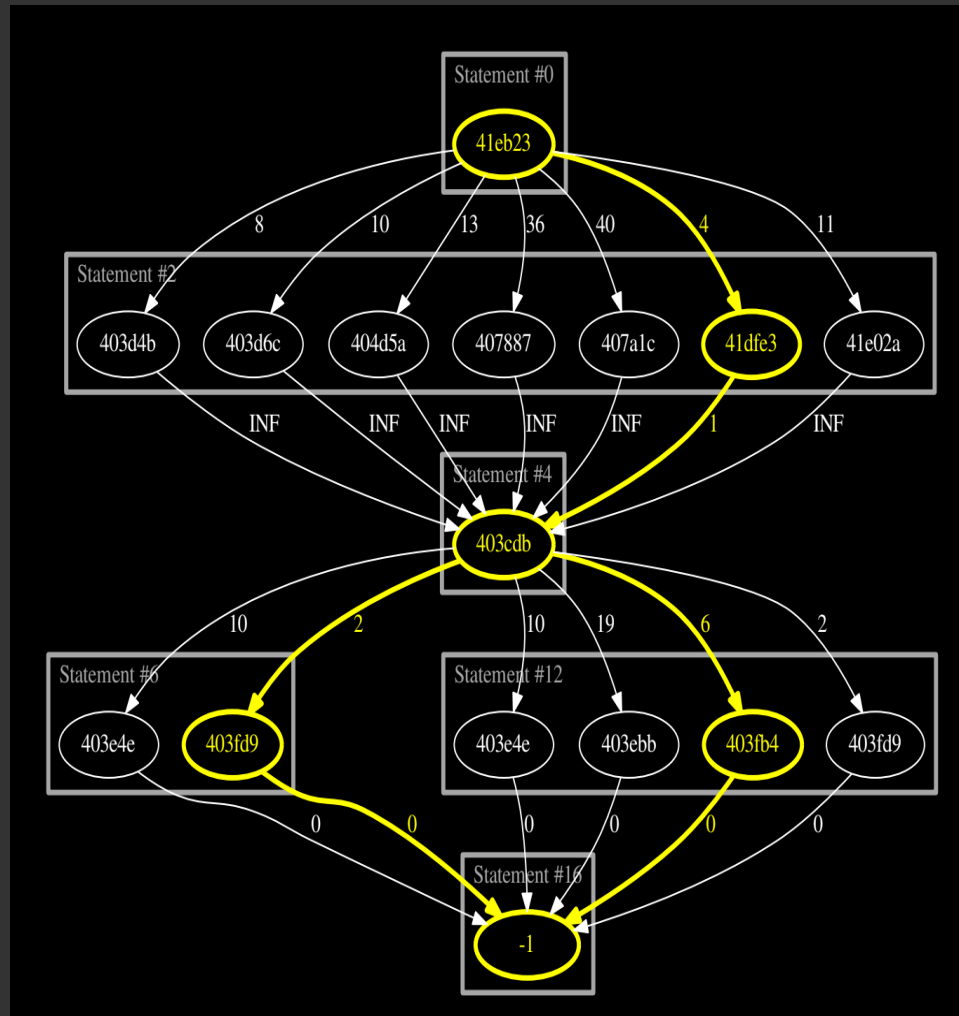


# Case study: inf loop on nginx

```
ngx_signal_handler()  
41C765: signals.signo == 0  
40E10F: ngx_time_lock != 0  
41C7B1: ngx_process 3 > 1  
41C9AC: ngx_cycle = $alloc_1  
         $alloc_1 > log = $alloc_2  
         $alloc_2 > log_level <= 5  
41CA18: signo == 17  
41CA4B: waitpid() return value != {0, 1}  
41cA50: ngx_last_process == 0  
41CB50: *($stack 0x03C) & 0x7F != 0  
41CB5B: $alloc_2 > log_level <= 1  
41CBE6: *($stack 0x03C + 1) != 2  
41CC48: ngx_accept_mutex_ptr == 0  
41CC5F: ngx_cycle > shared_memory.part.elts = 0  
         __r0 = r14 = 0  
41CC79: ngx_cycle > shared_memory.part.nelts <= 0  
41CC7F: ngx_cycle > shared_memory.part.next == 0
```



# Case study: if-else in nginx

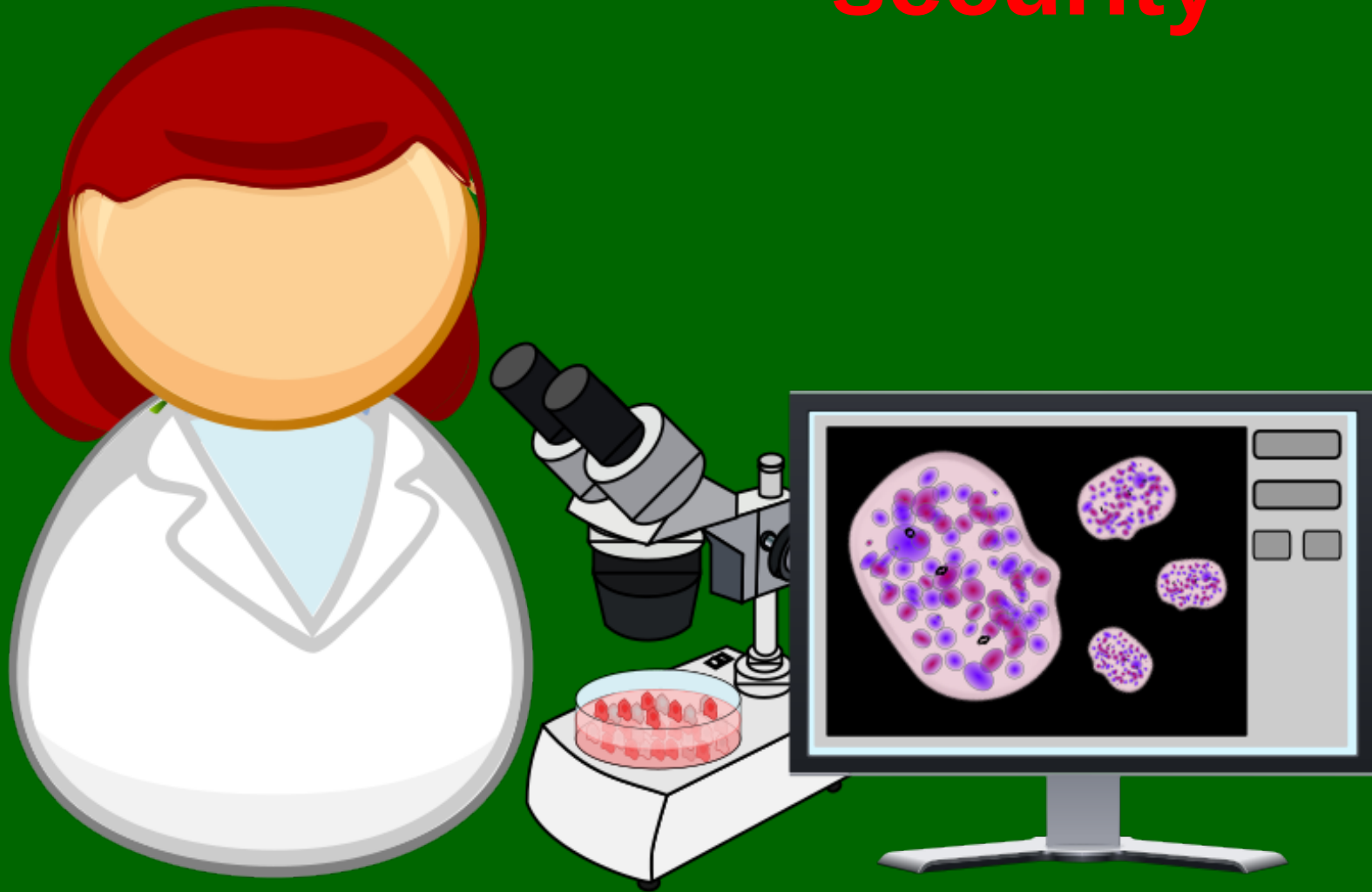


# BOP summary

- Block Oriented Programming
  - Automates Data-Only attacks
  - SPL: A language to express exploit payloads
  - Concolic execution algorithm stitches BOP gadgets
- We build exploits for 81% of the case studies
- Open source implementation (~14,000 LoC)

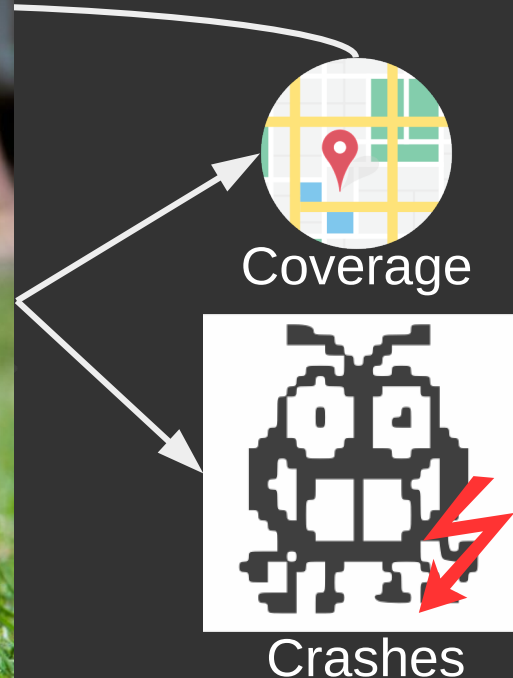
# Software testing: discover bugs

**security**



# Fuzz testing

- A random testing technique that mutates input to improve test coverage



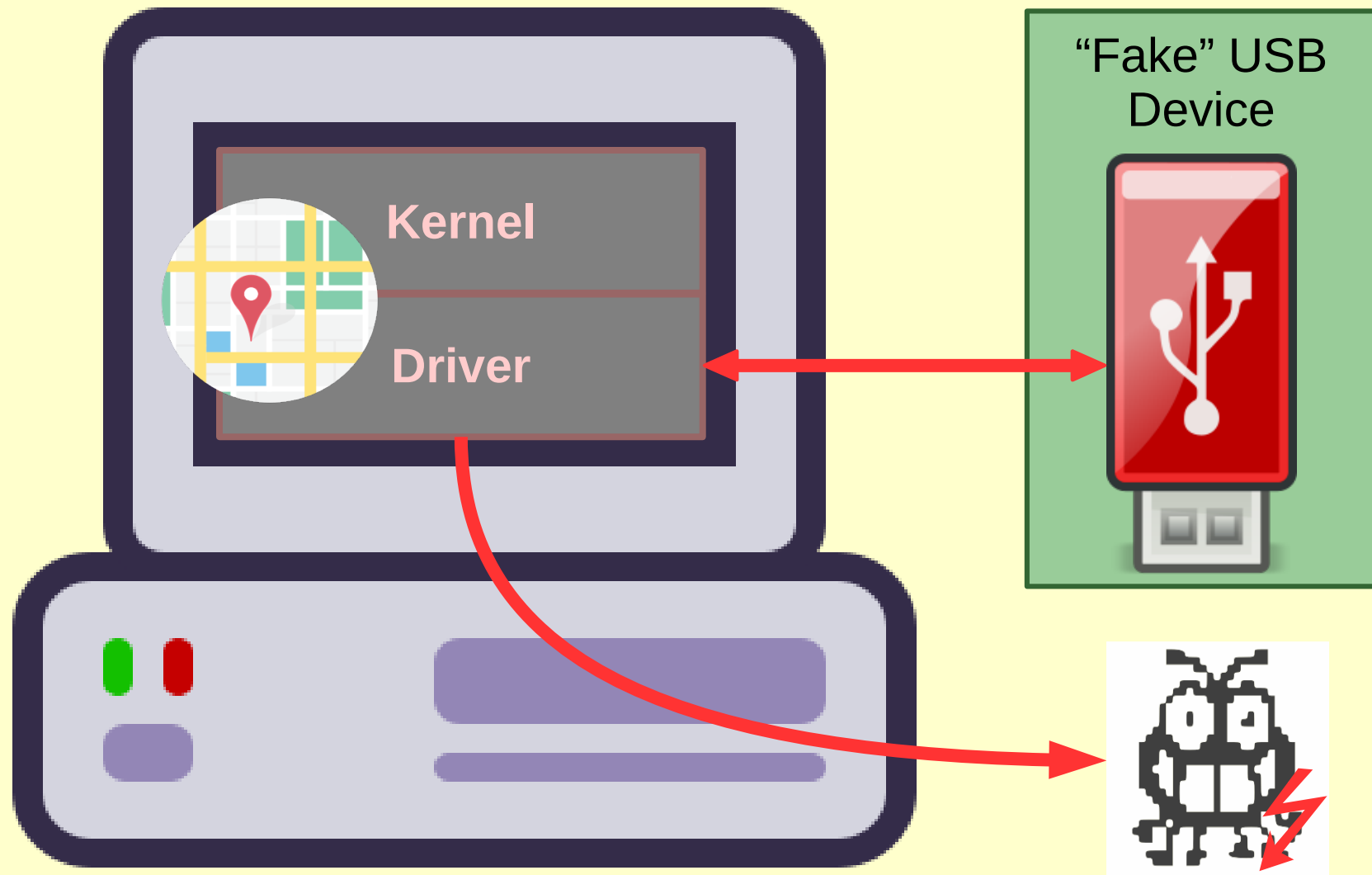
- State-of-art fuzzers use coverage as feedback to evolutionarily mutate the input

# Academic fuzzing research



# USB fuzz: explore peripheral space

## Virtual Environment



# USBFuzz Evaluation

- ~60 new bugs discovered in recent kernels
- 36 memory bugs (UaF / BoF)
- ~12 bugs fixed (with 9 CVEs)
- Bug reporting in progress

Type	Bug Info	#
Memory Bugs (36)	double-free	2
	NULL pointer dereference	8
	general protection	6
	slab-out-of-bounds access	6
	user-after-free access	16
Unexpected state reached (17)	INFO	6
	WARNING	9
	BUG	2



# Security testing hard-to-reach code

- Fuzzing is an effective way to automatically test programs for security violations (crashes)
  - Key idea: optimize for throughput
  - Coverage guides mutation
- BOP: assess exploitability
- USBFuzz: fuzz peripherals



<https://hexhive.epfl.ch>

<https://github.com/HexHive>

EPFL



# Vulnerable apps

Program	Vulnerability	Nodes	RegSet	RegMod	MemRd	MemWr	Call	Cond	Total
ProFTPD	CVE-2006-5815	27,087	40,143	387	1,592	199	77	3,029	45,427
nginx	CVE-2013-2028	24,169	31,497	1,168	1,522	279	35	3375	37,876
sudo	CVE-2012-0809	3,399	5,162	26	157	18	45	307	5715
orzhttpd	BID 41956	1,345	2,317	9	39	8	11	89	2473
wuftp	CVE-2000-0573	8,899	14,101	62	274	11	94	921	15,463
nullhttpd	CVE-2002-1496	1,488	2,327	77	54	7	19	125	2,609
opensshd	CVE-2001-0144	6,688	8,800	98	214	19	63	558	9,752
wireshark	CVE-2014-2299	74,186	124,053	639	1,736	193	100	4555	131276
apache	CVE-2006-3747	18,790	33,615	212	490	66	127	1,768	36,278
smbclient	CVE-2009-1886	166,081	265,980	1,481	6,791	951	119	28,705	304,027

**RegSet:** Register Assignment Gadgets  
**RegMod:** Register Modification Gadgets  
**MemRd:** Memory Read Gadgets  
**MemWr:** Memory Write Gadgets  
**Call:** Function/System Call Gadgets  
**Cond:** Conditional Statement Gadgets  
**Total:** Total number of Functional Gadgets

# SPL payloads

Payload	Description
<i>regset4</i>	Initialize 4 registers with arbitrary values
<i>regref4</i>	Initialize 4 registers with pointers to arbitrary memory
<i>regset5</i>	Initialize 5 registers with arbitrary values
<i>regref5</i>	Initialize 5 registers with pointers to arbitrary memory
<i>regmod</i>	Initialize a register with an arbitrary value and modify it
<i>memrd</i>	Read from arbitrary memory
<i>memwr</i>	Write to arbitrary memory
<i>print</i>	Display a message to stdout using write
<i>execve</i>	Spawn a shell through execve
<i>abloop</i>	Perform an arbitrarily long bounded loop utilizing regmod
<i>inloop</i>	Perform an infinite loop that sets a register in its body
<i>ifelse</i>	An if-else condition based on a register comparison
<i>loop</i>	Conditional loop with register modification